L Number	Hits	Search Text	DB	Time stamp
-	1501	707/2.ccls.	USPAT;	2003/11/20 13: 25
			US-PGPUB;	
			EPO; JPO;	
			IBM_TDB	
-	37	707/2.ccls. and (query adj cost)	USPAT;	2003/11/19 10:02
			US-PGPUB;	
1			ЕРО; ЛРО;	
			IBM_TDB	
-	34	(707/2.ccls. and (query adj cost)) and (@rlad <=20010228 @ad<=20010228)	USPAT;	2003/11/21 12:48
			US-PGPUB;	
1			EPO; JPO;	
			IBM_TDB	
-	78	707/2.ccls. and (generat\$3 adj index) and (@rlad <=20010228 @ad<=20010228)	USPĀT;	2003/11/19 10:09
			US-PGPUB:	
			EPO; JPO;	
			IBM_TDB	
-	50	(707/2.ccls. and (generat\$3 adj index) and (@rlad <=20010228 @ad<=20010228))	USPAT;	2003/11/19 10: 18
		and cost	US-PGPUB:	
			EPO; JPO;	
			IBM_TDB	
-	1	09/911,784	USPĀT:	2003/11/20 13: 24
			US-PGPUB:	
			ЕРО; ЈРО;	
			IBM_TDB	
-	1	6356889.URPN.	USPAT	2003/11/19 12:49
-	9	["5276870" "5822747" "5822750" "5864840" "5864842" "5930785" "6122644"	USPAT	2003/11/19 12:52
	•	"6199063" "6240406").PN.		
-	8	09/139843	USPAT;	2003/11/19 13: 31
			US-PGPUB:	
			ЕРО; ЈРО;	
			IBM_TDB	
-	6	09/087617	USPAT:	2003/11/19 13: 31
	_		US-PGPUB;	
			ЕРО; ЈРО;	
			IBM_TDB	
_	15	"5950186"	USPAT;	2003/11/19 15:48
			US-PGPUB;	2003, 22, 2, 25. 10
1			EPO; JPO;	
1			IBM_TDB	
-	14	"5960423"	USPAT:	2003/11/19 15: 51
			US-PGPUB;	2003/11/1/13.31
			EPO; JPO;	
			IBM_TDB	
1-	11	"5913207"	USPAT:	2003/11/19 15: 51
1			US-PGPUB;	2003/22/2/23:32
1			EPO; JPO;	
1			IBM TDB	
-	15	707/2.ccls. and (dynamic adj index\$3)	USPAT:	2003/11/20 08: 54
	•		US-PGPUB;	
1			EPO; JPO;	
			IBM_TDB	
1-	35	707/2,3,4,5.ccls. and (dynamic adj index\$3)	USPAT:	2003/11/20 08: 51
		· · · · - / - / · · · · · · · · · · · ·	US-PGPUB;	
1			EPO; JPO;	
			IBM_TDB	
-	32	(707/2,3,4,5.ccls. and (dynamic adj index\$3)) and (@rlad <=20010228	USPAT;	2003/11/20 08: 54
		@ad<=20010228)	US-PGPUB:	
		, O	EPO; JPO;	
			IBM_TDB	
[-	27	707/2.ccls. and (delet\$3 adj index\$3)	USPAT;	2003/11/20 08: 54
			US-PGPUB;	
			EPO; JPO;	
			IBM_TDB	
-	27	(707/2.ccls. and (delet\$3 adj index\$3)) and (@rlad <=20010228 @ad<=20010228)	USPAT;	2003/11/20 12: 11
			US-PGPUB;	
			EPO; JPO;	
			IBM_TDB	
-	54	5404510.URPN.	USPAT	2003/11/20 09: 06
-	80	(index adj maintenance) and (@rlad <=20010228 @ad<=20010228)	USPAT:	2003/11/20 05: 08
		,	US-PGPUB:	
			EPO; JPO;	
			IBM_TDB	
_	39	((index adj maintenance) and (@rlad <=20010228 @ad<=20010228)) and cost	USPAT;	2003/11/20 12:11
	27	mirror and interiorisation and limited TONIATED (Man TONIATED)) and COST	US-PGPUB;	2003/11/2012:11
			EPO; JPO;	
			IBM_TDB	
L			ם מודואומו	

555 combin\$3 adj index						
10 707/2.ccls. and (combinis) adj index)	-	1	6182079.URPN.	7.7	USPAT	2003/11/20 12:19
- 10 7072.ccls. and (combin53 adj index)	-	556	combin\$3 adj index		USPAT;	2003/11/20 13:24
10 767/2.ccls. and (combints) adj index 18M TDB					US-PGPUB;	
10 767/2.ccls. and (combints) adj index 18M TDB					ЕРО: ЛРО:	
10						
September Sept	-	10	707/2 ccls, and (combin\$3 adj index)			2003/11/20 15: 32
Second adj index Second adj	j	1	Torrando and too money			
18M TDB 18PA TDB						}
30 wider adj index SPAT	l	1				
US-PGPUB EPO, PPO, IBM, TDB USPAT US-PGPUB EPO, IPO, IBM, TDB US-PGPUB EPO, IPO, IPO, IBM, TDB US-PGPUB EPO, IPO, IPO, IPO, IPO, IPO, IPO, IPO, I	1.	30	wider adi index			2003/11/20 13:30
Solution		50	wider adj fridex			2003/11/2013.30
Separats						
5-5 generats3 adj second adj index		1				<u> </u>
S351 generat\$3 adj2 index US-PGPUB. EPO, IPO, IBM_TIDB USPAT. U	_	54	generat\$3 adi second adi indev			2003/11/20 13.41
For Pro Pro		, ,,,	generaliss and second and midex			2003/11/2013:41
S551 generat\$3 adj2 index	İ	İ				
S351 generats3 adj2 index					1''	
US-PGPUB, Property			concretes adis index			2002/11/20 12:41
FPO, IPO, IBM, TDB USPAT, USPOPUB, EPO, IPO, IBM, TDB USPAT, USPA	-	2221	generatiss adje index			2003/11/20 13:41
Minimal		1				
45 (generat53 adj second adj index) and (@rlad <=20010228 @ad<=20010228 USPAT, U		l			l —' — -'	
US-PGPUB; EPG; JPG; IBM, TIDB USPAT; U						
Combin\$3 adj (multiple plural\$4 two) adj index) and (@rlad <=20010228 USPAT, Qad<=20010228 USPAT, USP	-	45	(generat\$3 adj second adj index) and (@rlad <=20010228 @ad<=20010228)			2003/11/20 13:49
12						
12						
Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228) Combining adj indexes) and (@rlad <=20010228 @ad<=20010228)						
Pro, IPO, IBM, TDB USPAT	-	12	(combin\$3 adj (multiple plural\$4 two) adj index) and (@rlad <=20010228			2003/11/20 13:57
192 (index and database) ti. and (@rlad <=20010228 @ad<=20010228) IBM_TDB USPAT; USP-GPUB, EPO, IPO, IBM_TDB USPAT; USPAGPUB, EPO, IPO, IBM_TDB USPAT; USP-GPUB, EPO, IPO, IBM_TDB USP-GP			@ad<=20010228)			1
192 (index and database) ti. and (@rlad <=20010228 @ad<=20010228) IBM_TDB USPAT; USP-GPUB, EPO, IPO, IBM_TDB USPAT; USPAGPUB, EPO, IPO, IBM_TDB USPAT; USP-GPUB, EPO, IPO, IBM_TDB USP-GP					EPO; JPO;	
1	1	1				
US-PGPUB; EPO; PPO; IBM_TDB USPAT; USPAT; US-PGPUB; EPO; PPO; IBM_TDB USPAT; US	-	92	(index and database).ti. and (@rlad <=20010228 @ad<=20010228)		USPAT;	2003/11/20 14:55
1					US-PGPUB:	
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11 3 (*4805099**) **494510**).PN. USPAT US	1	ļ				
10221 Sys0a23.URPN. (index and reuse\$1) and (@rlad <=20010228 @ad<=20010228) USPAT; US-PGPUB; EPO, IPO; IBM_TDB USPAT; US-PGPUB; EPO, IPO; IBM_TDB USPAT; US-PGPUB; EPO, IPO; IBM_TDB USPAT; US-PGPUB; EPO, IPO; IBM_TDB USPAT; US-PGPUB; EPO; IPO; IBM_TDB USPAT; US-PGPUB; IBM_TDB USPAT; US-PGPUB; IBM_TDB USPAT; US-PGPUB; IBM_TDB US-PAT; IBM_TDB US-PAT; IBM_TDB US-PAT	-	3	("4805099" "4956774" "5404510").PN.			2003/11/20 14: 04
10221		11				
US-PGPUB; EPO; IPO; IBM_TDB USPAT; US-PGPUB; EPO; IPO; IPO; IBM_TDB USPAT; US-PGPUB; EPO; IPO; IPO; IPO; IBM_TDB U	-	1				
EPO; IPO; IBM TDB USPAT; USPGPUB; EPO; IPO; IBM TDB USPAT; USPAT; USPAT; USPAT; USPAT; USPAT; USPAT; USPGPUB; EPO; IPO; IBM TDB USPAT; USPAT; USPGPUB; EPO; IPO; IBM TDB USPAT; USPGPUB; EPO; IP			,, (O)			
29 (index adj reuse\$1) and (@rlad <=20010228 @ad<=20010228)	}					
- 135 (index adj reuse\$1) and (@rlad <=20010228 @ad<-20010228) USPAT, US-PGPUB, EPO, JPO, IBM_TDB USPAT (index near reuse\$1) and (@rlad <=20010228 @ad<-20010228) USPAT, US-PGPUB, EPO, JPO, IBM_TDB USPAT, US-PGPUB, EPO, JPO, IB						
US-PGPUB; EPO; PO; IBM_TDB USPAT; US-PGPUB; EPO; PO; IBM_TDB US-PAT; US-PGPUB; EPO; PO; PO; PO; PO; PO; PO; PO; PO; PO;	<u>-</u>	29	(index adi reuse\$1) and (@rlad <=20010228 @ad<=20010228)			2003/11/20 15:00
Fep Fo Fo Fo Fo Fo Fo Fo F		1	Titalit day readers, died (with the terminal ter			2003/22/20 25:00
- \$4 \$404510.URPN. (index near reuse\$1) and (@rlad <=20010228 @ad<=20010228) USPAT; US						
- \$4 \$4 \$4 \$4 \$4 \$54 \$40 \$4 \$10 \$4 \$10 \$4 \$10 \$10 \$28 \$(a) \$4 \$4 \$10 \$10 \$10 \$10 \$10 \$10 \$10 \$10 \$10 \$10						
- 135 (index near reuse\$1) and (@rlad <=20010228 @ad<=20010228)	_	54	5404510 LTRPN			2003/11/20 14:58
US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO;	_				_	1
- 135 707/2.ccls. and (index with range) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 133 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 14 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228) - 15 (index same (retrieval adj condition) same combin\$3) and (@rlad <=20010228) - 16 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 17 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 18 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 19 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 11 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 12 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 2003/11/20 16: 12 - 2003/11/20 16: 12 - 2003/11/21 08: 16 - 2003/11/21 08: 28 - 2003/11/21 08: 28 - 2003/11/21 09: 32			,			
- 135 707/2.ccls. and (index with range) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 14 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 @ad<=20010228) - 15 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 @ad<=20010228) - 16 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 17 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 18 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 19 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 10 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 11 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 12 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 13 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 @ad<=20010228) - 14 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 15 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 @ad<=20010228) - 16 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 17 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 18 (index with (retrieval adj condition) with combin\$3 and (@rlad <=20010228 @ad<=20010228) - 17 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228)					,	
- 135 707/2.ccls. and (index with range) - 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228) - 8 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228	1					
- 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IPO; IPO; IPO; IPO; IPO; IPO; IPO; I	_	135	707/2 ccls, and (index with range)			2003/11/20 15, 33
EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO;		***	,			2003/11/2013:33
- 131 (707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228) - 8 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 USPAT; US-PGPUB; EPO; IPO; IBM_TDB USPAT; @ad<=20010228) - 11 (index same (retrieval adj condition) same combin\$3) and (@rlad <=20010228 USPAT; US-PGPUB; EPO; IPO; IBM_TDB USPAT; US-PGPUB; EPO; IPO; IPO; IBM_TDB USPAT; US-PGPUB; EPO; IPO; IPO; IPO; IPO; IPO; IPO; IPO; I						
- 40 707/2.ccls. and (index with range)) and (@rlad <=20010228 @ad<=20010228) - 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228) - 8 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; ISM_TDB USPAT; US-PGPUB; ISM_TDB USPAT; U						
- 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228)	_	121	1707/2 ccls and (index with range)) and (Orlad - 20010229 Oad - 20010229)			2003/11/20 14:04
- 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228) - 8 (index with (retrieval adj condition) with combin\$3) and (@rlad <=20010228 USPAT; US-PGPUB; EPO, JPO; IBM_TDB USPAT; US-PGPUB; EPO, JPO;		131	1, 2, 2, 2, 2013. and finder with range// and florial <=20010226 (mau<=20010228)	'		2003/11/20 10:04
- 40 707/2,3,4.ccls. and (range adj predicate) and (@rlad <=20010228 @ad<=20010228)						
- 4 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 4 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 4 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 5 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 6 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 7 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228)						
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Beautiful Combiner Combiner	~	40	101/2,3,4.0015. and frange and predicate) and (@frad <=20010228 @ad<=2001	0228)		2003/11/20 16:11
BM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB		[
Solution Combining adj indexes Combining adj indexes Combining and (Combining adj indexes) Combining adj indexes Combining and (Combining adj indexes) Combining and (Combining and (Combin						
Combining adj indexes) and (@rlad <=20010228) US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; ISPAT; US-PGPUB; EPO; ISPAT; US-PGPUB; EPO; ISPAT;		1	(index with (matrice) and and conditions with a well-tops and (O.1.3.)			2002.11.22.11
- 11 (index same (retrieval adj condition) same combin\$3) and (@rlad <=20010228	-	8			USPA1;	2003/11/20 16: 12
- (index same (retrieval adj condition) same combin\$3) and (@rlad <=20010228			(@ad<=20010228)			
- (index same (retrieval adj condition) same combin\$3) and (@rlad <=20010228						
- 4 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 4 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) - 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO;						
- 4 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO;	-	11		8		2003/11/21 08:16
- 4 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; IBM_TOB USPAT; US-PGPUB; EPO; JPO; ISPO;			(@ad<=20010228)			
- 4 (index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228) - 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) USPĀT; US-PGPUB; EPO; JPO; IBM_TDB USPĀT; US-PGPUB; EPO; JPO; PO; JPO;						
US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO; US-PGPUB; EPO; JPO;		1				
- 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) EPO; JPO; IBM_TDB USPAT; US-PGPUB; EPO; JPO;	-	4	(index adj selection adj algorithm) and (@rlad <=20010228 @ad<=20010228))		2003/11/21 08: 28
- 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) IBM_TDB USPAT; US-PGPUB; EPO; JPO;		1				
- 66 (combining adj indexes) and (@rlad <=20010228 @ad<=20010228) USPAT; US-PGPUB; EPO; JPO;		1			ЕРО; ЈРО;	
US-PGPUB; EPO; JPO;						
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			ЕРО; ЛРО;	
			IBM_TDB	
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-	3		USPAT	2003/11/21 10: 29
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			IBM_TDB	
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		"6006220").PN.		
-	15		USPAT	2003/11/21 13:54
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On the Selection of Secondary Indices in Relational Databases (1993) (Make Corrections) (6 citations)

Sunil Choenni, Henk M. Blanken, Thiel Chang Data Knowledge Engineering



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Abstract: An important problem in the physical design of databases is the selection of secondary indices. In general, this problem can not be solved in an optimal way due to the complexity of the selection process. Often use is made of heuristics such as the well-known ADD and DROP algorithms. In this paper it will be shown that frequently used cost functions can be classified as super- or submodular functions. For these functions several mathematical properties have been derived which reduce the... (Update)

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- 1.2: Index Configurations in Object-Oriented Databases Choenni, Bertino, Blanken.. (Correct)
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BibTeX entry: (Update)

```
@article{ choenni93selection,
   author = "Sunil Choenni and Henk M. Blanken and T. Chang",
   title = "On the Selection of Secondary Indices in Relational Databases",
   journal = "Data Knowledge Engineering",
   volume = "11",
   number = "3",
   pages = "207-",
   year = "1993",
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Index Selection for OLAP - Gupta, Harinarayan, Rajaraman, Ullman (1997) (Correct) (68 citations)

Index Selection for OLAP #Himanshu Gupta Venky Harinarayan

This paper is the #rst to explore the index-selection problem and automate it with provably Now consider integrating the subcube selection and index selection into one step. We use the www.cise.ufl.edu/~igreenbe/research/../papers/23.pdf

Cluster-based Language Models for Distributed Retrieval - Xu, Croft (1999) (Correct) (29 citations) data structure is called the collection selection index #23#Collection selection consists of in the same time period. The collection selection index summarizes each collection as a whole. In is illustrated by Figure 1. Collection Selection Index collection 1 collection 2 collection3 www.cs.virginia.edu/~cyberia/cs851/papers/xu-clstr.pdf

Cluster-based Language Models For Distributed Retrieval - Xu, Croft (1999) (Correct) (25 citations) data structure is called the collection selection index [23]Collection selection consists of in the same time period. The collection selection index summarizes each collection as a whole. In is illustrated by Figure 1. Collection Selection Index collection 1 collection 2 collection3 ciir.cs.umass.edu/pubfiles/ir-153.ps

Comparing the Performance of Database Selection.. - French, Powell, Callan, .. (1999) (Correct) (15 citations) search, the CORI approach creates a database selection index in which each database is represented by its www.cs.virginia.edu/~cyberia/papers/SIGIR99.ps

Subtopic Structuring for Full-Length Document Access - Hearst, Plaunt (1993) (Correct) (55 citations) can im- prove on categorization and index selection tasks. We then describe an experiment in www.ai.mit.edu/people/jimmylin/papers/Hearst93.pdf

Optimization of Run-time Management of Data Intensive Web .. - Florescu, Levy, Suciu. (1999) (Correct) (12 citations)

as view materialization [24, 13, 12, 14, 6] index selection, data caching [16, 15, 8] multiple query 14. 6]multiple query optimization [23] and index selection. All these techniques are of course www.research.att.com/~suciu/strudel/external/files/ F269651526.ps

Optimization of Run-time Management of Data Intensive.. - Florescu, Levy, Suciu.. (1999) (Correct) (12 citations) view materialization [13, 23, 10, 12, 11, 6]index selection [7]function caching [16, 14]multiple 12, 11, 6]multiple query optimization [22]index selection [7]and parameterized query optimization rodin.inria.fr/dataFiles/FLSY99b.ps

Effective Retrieval with Distributed Collections - Xu, Callan (1998) (Correct) (19 citations) is to use phrase information in the collection selection index and the other is query expansion. Both 8/98 \$5.00. creates a collection selection index. The collection selection index consists of a collection selection index. The collection selection index consists of a set of virtual documents. www.cs.umass.edu/~xu/sigir98-final.ps

Physical Database Design for Data Warehouses - Labio, Quass, Adelberg (1997) (Correct) (17 citations) the total down time. We call this the view index selection (VIS) problem. We present an exhaustive materialized views, view maintenance, index selection, and physical database design. 1 views it is necessary to consider the view selection and index selection together. If view selection is ftp.dblab.ntua.gr/pub/dwg/view.ps.gz

The Dimensional Fact Model: A Conceptual Model For Data.. - Golfarelli, Maio, Rizzi (1998) (Correct) (10 citations) such as materialization of views 2,15 and index selection 13,16 no significant effort has been 13. H. Gupta, V. Harinarayan and A. Rajaraman, Index selection for OLAP, Proc. Int. Conf. Data Engineering,

ftp-db.deis.unibo.it/pub/stefano/ijcis98.ps

A Comprehensive Approach to Horizontal Class Fragmentation in .. - Ezeife, Barker (1995) (Correct) (18 citations) of data. The work of Bertino and Kim [1] on index selection is also complementary to our work, but www.cs.uwindsor.ca/users/c/cezeife/journal1.ps

Rethinking Database System Architecture: Towards a... - Chaudhuri, Weikum (2000) (Correct) (3 citations) of the application. These knobs include **index selection**, data placement across parallel disks, and www-dbs.cs.uni-sb.de/public_html/papers/cairo-final.pdf

A Methodological Framework for Data Warehouse Design - Golfarelli, Rizzi (1998) (Correct) (7 citations) [1] 8]materialization of views [2] 9] and index selection [7] 10]no significant effort has been provided by the DBMS to be taken into account. Index selection has a crucial role in determining the DW ftp-db.deis.unibo.it/pub/stefano/dolap98.ps.gz

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Relevance scale

An efficient expected cost algorithm for dynamic indexing of spatial objects Gopi K. Attaluri

October 1994 Proceedings of the 1994 conference of the Centre for Advanced Studies on Collaborative research

Full text available: pdf(173.67 KB) Additional Information: full citation, abstract, references, index terms

Spatial object management is important in numerous application areas, including geographic data, CAD, and VLSI design. It often involves insertion, deletion, and range search of spatial objects, so requires efficient dynamic indexing of objects. This paper describes an (in memory) indexing algorithm developed in the context of concurrency control for large, unstructured objects in non-traditional database applications. It supports insertion, deletion, and range search of orthogonal spatial object ...

Keywords: concurrency control, indexing, range locking, runtime complexity, spatial databases

2 XML indexing and compression: ViST: a dynamic index method for querying XML data by tree structures



Haixun Wang, Sanghyun Park, Wei Fan, Philip S. Yu

June 2003 Proceedings of the 2003 ACM SIGMOD international conference on on Management of data

Full text available: pdf(244.47 KB) Additional Information: full citation, abstract, references, index terms

With the growing importance of XML in data exchange, much research has been done in providing flexible query facilities to extract data from structured XML documents. In this paper, we propose ViST, a novel index structure for searching XML documents. By representing both XML documents and XML queries in structure-encoded sequences, we show that querying XML data is equivalent to finding subsequence matches. Unlike index methods that disassemble a query into multiple sub-queries, and then *joi* ...

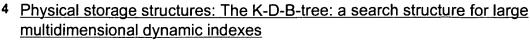
3 Segment indexes: dynamic indexing techniques for multi-dimensional interval data Curtis P. Kolovson, Michael Stonebraker



April 1991 ACM SIGMOD Record, Proceedings of the 1991 ACM SIGMOD international conference on Management of data, Volume 20 Issue 2

Full text available: pdf(1.02 MB)

Additional Information: full citation, references, citings, index terms





John T. Robinson

April 1981 Proceedings of the 1981 ACM SIGMOD international conference on Management of data

Full text available: 7 pdf(723.91 KB) Additional Information: full citation, abstract, references, citings

The problem of retrieving multikey records via range queries from a large, dynamic index is considered. By *large* it is meant that most of the index must be stored on secondary memory. By *dynamic* it is meant that insertions and deletions are intermixed with queries, so that the index cannot be built beforehand. A new data structure, the *K-D-B-tree*, is presented as a solution to this problem. K-D-B-trees combine properties of K-D-trees and B-trees. It is expected that the mult ...

5 Physical database design: R-trees: a dynamic index structure for spatial searching. Antonin Guttman



June 1984 Proceedings of the 1984 ACM SIGMOD international conference on Management of data

Full text available: pdf(1.07 MB)

Additional Information: full citation, abstract, references, citings

In order to handle spatial data efficiently, as required in computer aided design and geodata applications, a database system needs an index mechanism that will help it retrieve data items quickly according to their spatial locations However, traditional indexing methods are not well suited to data objects of non-zero size located m multi-dimensional spaces In this paper we describe a dynamic index structure called an R-tree which meets this need, and give algorithms for searching and updating ...

6 Comparison of access methods for time-evolving data

Betty Salzberg, Vassilis J. Tsotras

June 1999 ACM Computing Surveys (CSUR), Volume 31 Issue 2

Full text available: pdf(529.53 KB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> <u>terms</u>

This paper compares different indexing techniques proposed for supporting efficient access to temporal data. The comparison is based on a collection of important performance criteria, including the space consumed, update processing, and query time for representative queries. The comparison is based on worst-case analysis, hence no assumptions on data distribution or query frequencies are made. When a number of methods have the same asymptotic worst-case behavior, features in the methods tha ...

Keywords: I/O performance, access methods, structures, temporal databases

7 Multidimensional access methods

Volker Gaede, Oliver Günther

June 1998 ACM Computing Surveys (CSUR), Volume 30 Issue 2

Full text available: pdf(1.05 MB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> <u>terms</u>

Search operations in databases require special support at the physical level. This is true for conventional databases as well as spatial databases, where typical search operations include the point query (find all objects that contain a given search point) and the region query (find all objects that overlap a given search region). More than ten years of spatial database research have resulted in a great variety of multidimensional access methods to support ...

Keywords: data structures, multidimensional access methods

8	The SR-tree: an index structure for high-dimensional nearest neighbor queries	_
	Norio Katayama, Shin'ichi Satoh	
	June 1997 ACM SIGMOD Record , Proceedings of the 1997 ACM SIGMOD international	
	conference on Management of data, Volume 26 Issue 2 Full text evaluable: To adf(4.44 MR) Additional Information: full citation, abstract, references, citings, index	
	Full text available: pdf(1.41 MB) Additional information. information, abstract, references, citings, index terms	
	Recently, similarity queries on feature vectors have been widely used to perform content-based retrieval of images. To apply this technique to large databases, it is required to develop multidimensional index structures supporting nearest neighbor queries efficiently. The SS-tree had been proposed for this purpose and is known to outperform other index structures such as the R*-tree and the K-D-B-tree. One of its most important features is that it employs bounding spheres rather than boundi	
9	On packing R-trees	_
	Ibrahim Kamel, Christos Faloutsos	
	December 1993 Proceedings of the second international conference on Information and	
	knowledge management Full text available: pdf(929.84 KB) Additional Information: full citation, references, citings, index terms	
	r dir toxt d'aliable. parteze.e 1718/	
10	Parallel R-trees	_
	Ibrahim Kamel, Christos Faloutsos	-
	June 1992 ACM SIGMOD Record, Proceedings of the 1992 ACM SIGMOD international conference on Management of data, Volume 21 Issue 2	
	Full text available: pdf(991.13 KB) Additional Information: full citation, abstract, references, citings, index	
	terms	
	We consider the problem of exploiting parallelism to accelerate the performance of spacial access methods and specifically, R-trees [11]. Our goal is to design a server for spatial data, so that to maximize the throughput of range queries. This can be achieved by (a) maximizing parallelism for large range queries, and (b) by engaging as few disks as possible on point queries [22]. We propose a simple hardware architecture consisting of one processor with several disks attached to	
11	High performance clustering based on the similarity join	_
	Christian Böhm, Bernhard Braunmüller, Markus Breunig, Hans-Peter Kriegel November 2000 Proceedings of the ninth international conference on Information and knowledge management	
	Full text available: pdf(134.57 KB) Additional Information: full citation, references, citings, index terms	
	Keywords : clustering, data mining, database primitives, multidimensional index structure, similarity join	

Full text available: pdf(588.72 KB) Additional Information: full citation, references, index terms

12 <u>Searching spatial objects with index by dimensional projections</u>
Xiaoming Cheng, Huizhu Lu, G. E. Hedrick

April 1992 Proceedings of the 1992 ACM/SIGAPP Symposium on Applied computing: technological challenges of the 1990's

13 Special issue on spatial database systems: An introduction to spatial database	
<u>systems</u>	
Ralf Hartmut Güting October 1994 The VLDB Journal — The International Journal on Very Large Data Bases, Volume 3 Issue 4	
Full text available: pdf(2.50 MB) Additional Information: full citation, abstract, references, citings	
We propose a definition of a spatial database system as a database system that offers spatial data types in its data model and query language, and supports spatial data types in its implementation, providing at least spatial indexing and spatial join methods. Spatial database systems offer the underlying database technology for geographic information systems and other applications. We survey data modeling, querying, data structures and algorithms, and system architecture for such systems. The em	
14 Redundancy in spatial databases J. A. Orenstein	
June 1989 ACM SIGMOD Record, Proceedings of the 1989 ACM SIGMOD international conference on Management of data, Volume 18 Issue 2	
Full text available: pdf(1.37 MB) Additional Information: full citation, abstract, references, citings, index terms	
Spatial objects other than points and boxes can be stored in spatial indexes, but the techniques usually require the use of approximations that can be arbitrarily bad. This leads to poor performance and highly inaccurate responses to spatial queries. The situation can be improved by storing some objects in the index redundantly. Most spatial indexes permit no flexibility in adjusting the amount of redundancy. Spatial indexes based on z-order permit this flexibility. Accuracy of the query re	
15 <u>Scalable integrated region-based image retrieval using IRM and statistical clustering</u> James Z. Wang, Yanping Du January 2001 Proceedings of the first ACM/IEEE-CS joint conference on Digital libraries	
Full text available: pdf(1.73 MB) Additional Information: full citation, abstract, references, citings, index terms	
Statistical clustering is critical in designing scalable image retriev al systems. In this paper, we present a scalable algorithm for indexing and retrieving images based on region segmentation. The method uses statistical clustering on region features and IRM (Integrated Region Matching), a measure developed to evaluate overall similarity between images that incorporates properties of all the regions in the images by a region-matching scheme. Compared with retrieval based on individual	
Keywords : clustering, content-based image retrieval, integrated region matching, segmentaton, wavelets	
16 On effective multi-dimensional indexing for strings H. V. Jagadish, Nick Koudas, Divesh Srivastava May 2000 ACM SIGMOD Record, Proceedings of the 2000 ACM SIGMOD international conference on Management of data, Volume 29 Issue 2	
Full text available: pdf(1.15 MB) Additional Information: full citation, abstract, references, citings, index terms	
As databases have expanded in scope from storing purely business data to include XML documents, product catalogs, e-mail messages, and directory data, it has become increasingly important to search databases based on wild-card string matching: prefix matching, for example, is more common (and useful) than exact matching, for such data. In many cases, matches need to be on multiple attributes/dimensions, with correlations	

17 Simple QSF-trees: an efficient and scalable spatial access method Byunggu Yu, Ratko Orlandic, Martha Evens November 1999 Proceedings of the eighth international conference on Information and knowledge management Additional Information: full citation, abstract, references, citings, index Full text available: pdf(1.33 MB) terms The development of high-performance spatial access methods that can support complex operations of large spatial databases continues to attract considerable attention. This paper introduces QSF-trees, an efficient and scalable structure for indexing spatial objects, which has some important advantages over R*-trees. QSF-trees eliminate overlapping of index regions without forcing object clipping or sacrificing the select ... Keywords: database management, point access methods, spatial access methods, spatial database, topological relations 18 Distance browsing in spatial databases Gísli R. Hjaltason, Hanan Samet June 1999 ACM Transactions on Database Systems (TODS), Volume 24 Issue 2 Additional Information: full citation, abstract, references, citings, index Full text available: pdf(460.81 KB) terms We compare two different techniques for browsing through a collection of spatial objects stored in an R-tree spatial data structure on the basis of their distances from an arbitrary spatial query object. The conventional approach is one that makes use of a k-nearest neighbor algorithm where k is known prior to the invocation of the algorithm. Thus if m < kneighbors are needed, the k-nearest neighbor alg ... Keywords: R-trees, distance browsing, hiearchical spatial data structures, nearest neighbors, ranking 19 Indexing large metric spaces for similarity search queries Tolga Bozkaya, Meral Ozsoyoglu September 1999 ACM Transactions on Database Systems (TODS), Volume 24 Issue 3 Additional Information: full citation, abstract, references, citings, index Full text available: pdf(281.78 KB) terms, review One of the common queries in many database applications is finding approximate matches to a given query item from a collection of data items. For example, given an image database, one may want to retrieve all images that are similar to a given query image. Distance-based index structures are proposed for applications where the distance computations between objects of the data domain are expensive (such as high-dimensional data) and the distance function is metric. In this paper we consider ... ²⁰ Indexing schemes for random points Elias Koutsoupias, David Taylor January 1999 Proceedings of the tenth annual ACM-SIAM symposium on Discrete algorithms Full text available: 🔂 pdf(838.33 KB) Additional Information: full citation, references, citings, index terms

between the dimensions. Traditional multi-dimensional index structures, ...

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Relevance scale

1 A cost model for query processing in high dimensional data spaces
Christian Böhm

June 2000 ACM Transactions on Database Systems (TODS), Volume 25 Issue 2

Full text available: pdf(362.22 KB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> <u>terms</u>, <u>review</u>

During the last decade, multimedia databases have become increasingly important in many application areas such as medicine, CAD, geography, and molecular biology. An important research topic in multimedia databases is similarity search in large data sets. Most current approaches that address similarity search use the feature approach, which transforms important properties of the stored objects into points of a high-dimensional space (feature vectors). Thus, similarity search is transformed ...

Keywords: cost model, multidimensional index

² Minimum cost selection of secondary indexes for formatted files

Henry D. Anderson, P. Bruce Berra

Management of data

March 1977 ACM Transactions on Database Systems (TODS), Volume 2 Issue 1

Full text available: pdf(1.74 MB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> terms

Secondary indexes are often used in database management systems for secondary key retrieval. Although their use can improve retrieval time significantly, the cost of index maintenance and storage increases the overhead of the file processing application. The optimal set of indexed secondary keys for a particular application depends on a number of application dependent factors. In this paper a cost function is developed for the evaluation of candidate indexing choices and applied to the opti ...

Keywords: Boolean query, access methods, access path, cost function, data management, database, file design, file organization, inverted file, inverted index, maintenance, optimization, retrieval, secondary index, secondary key, secondary key access

Query processing: Factorizing complex predicates in queries to exploit indexes Surajit Chaudhuri, Prasanna Ganesan, Sunita Sarawagi June 2003 Proceedings of the 2003 ACM SIGMOD international conference on on



Full text available: pdf(240.56 KB) Additional Information: full citation, abstract, references, index terms

Decision-support applications generate queries with complex predicates. We show how the *factorization* of complex query expressions exposes significant opportunities for exploiting available indexes. We also present a novel idea of relaxing predicates in a complex condition to create possibilities for factoring. Our algorithms are designed for easy integration with existing query optimizers and support multiple optimization levels, providing different tradeoffs between plan complexity and ...

4 Query processing: A characterization of the sensitivity of query optimization to storage access cost parameters



Frederick R. Reiss, Tapas Kanungo

June 2003 Proceedings of the 2003 ACM SIGMOD international conference on on Management of data

Full text available: 🔂 pdf(255.35 KB) Additional Information: full citation, abstract, references, index terms

Most relational query optimizers make use of information about the costs of accessing tuples and data structures on various storage devices. This information can at times be off by several orders of magnitude due to human error in configuration setup, sudden changes in load, or hardware failure. In this paper, we attempt to answer the following questions: • Are inaccurate access cost estimates likely to cause a typical query optimizer to choose a suboptimal query plan? • If an optimizer ...

Keywords: autonomic computing, computational geometry, databases, parametric query optimization, storage systems

5 Location awareness and moving objects: Efficient placement of geographical data over broadcast channel for spatial range query under quadratic cost model Jianting Zhang, Le Gruenwald



September 2003 Proceedings of the 3rd ACM international workshop on Data engineering for wireless and mobile access

Full text available: pdf(326.37 KB) Additional Information: full citation, abstract, references, index terms

Data broadcasting is well known for its excellent scalability. Most geographical data, such as weather and traffic, is public information that has a large amount of potential users which makes it very suitable for broadcast. The query response time is greatly affected by the order in which data items are being broadcast. This paper proposes an efficient method to place geographical data items over broadcast channel that reduces access time for spatial range queries on them. This paper then perfo ...

Keywords: cost model, data broadcast, geographical information, mobile computing, optimization, query processing

⁶ A cost model for similarity queries in metric spaces

Paolo Ciaccia, Marco Patella, Pavel Zezula



Full text available: pdf(1.12 MB) Additional Information: full citation, references, citings, index terms

7 Improved query performance with variant indexes Patrick O'Neil, Dallan Quass

June 1997 ACM SIGMOD Record, Proceedings of the 1997 ACM SIGMOD international conference on Management of data, Volume 26 Issue 2

Full text available: pdf(1.54 MB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> terms

The read-mostly environment of data warehousing makes it possible to use more complex indexes to speed up queries than in situations where concurrent updates are present. The current paper presents a short review of current indexing technology, including row-set representation by Bitmaps, and then introduces two approaches we call Bit-Sliced indexing and Projection indexing. A Projection index materializes all values of a column in RID order, and a Bit-Sliced index essentially takes an orth ...

Indexing large metric spaces for similarity search queries

Tolga Bozkaya, Meral Ozsoyoglu

September 1999 ACM Transactions on Database Systems (TODS), Volume 24 Issue 3

Full text available: pdf(281.78 KB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> <u>terms</u>, <u>review</u>

One of the common queries in many database applications is finding approximate matches to a given query item from a collection of data items. For example, given an image database, one may want to retrieve all images that are similar to a given query image. Distance-based index structures are proposed for applications where the distance computations between objects of the data domain are expensive (such as high-dimensional data) and the distance function is metric. In this paper we consider ...

⁹ Efficiency: Compression of inverted indexes For fast query evaluation

Falk Scholer, Hugh E. Williams, John Yiannis, Justin Zobel

August 2002 Proceedings of the 25th annual international ACM SIGIR conference on Research and development in information retrieval

Full text available: pdf(174.13 KB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> terms

Compression reduces both the size of indexes and the time needed to evaluate queries. In this paper, we revisit the compression of inverted lists of document postings that store the position and frequency of indexed terms, considering two approaches to improving retrieval efficiency: better implementation and better choice of integer compression schemes. First, we propose several simple optimisations to well-known integer compression schemes, and show experimentally that these lead to significan ...

Keywords: index compression, integer coding, inverted indexes, retrieval efficiency

10 Document query processing strategies: cost evaluation and heuristics

E. Bertino, S. Gibbs, F. Rabitti

April 1988 Conference Sponsored by ACM SIGOIS and IEEECS TC-OA on Office information systems

Full text available: pdf(969.16 KB) Additional Information: full citation, abstract, references, index terms

This paper describes query processing strategies used in a system where queries are specified on formatted data and text components of documents. The system provides different access methods: indexes for formatted data and signature files for the text. Four basic document queries, representative of a wide range of queries, are examined in detail. Strategies for query processing are presented together with an evaluation of the costs of the various strategies. Finally, general heurist ...

11 The onion technique: indexing for linear optimization queries

Yuan-Chi Chang, Lawrence Bergman, Vittorio Castelli, Chung-Sheng Li, Ming-Ling Lo, John R. Smith

May 2000 ACM SIGMOD Record, Proceedings of the 2000 ACM SIGMOD international

conference on Management of data, Volume 29 Issue 2

Full text available: pdf(326.62 KB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> <u>terms</u>

This paper describes the Onion technique, a special indexing structure for linear optimization queries. Linear optimization queries ask for top-N records subject to the maximization or minimization of linearly weighted sum of record attribute values. Such query appears in many applications employing linear models and is an effective way to summarize representative cases, such as the top-50 ranked colleges. The Onion indexing is based on a geometric property of convex hull, which guarantees ...

Keywords: database indexing, linear optimization

12 Optimization of object-oriented recursive queries using cost-controlled strategies
Rosana S. G. Lanzelotte, Patrick Valduriez, Mohamed Zaït

June 1992 ACM SIGMOD Record, Proceedings of the 1992 ACM SIGMOD international conference on Management of data, Volume 21 Issue 2

Full text available: pdf(1.07 MB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> terms

Object-oriented data models are being extended with recursion to gain expressive power. This complicates the optimization problem which has to deal with recursive queries on complex objects. Because unary operations invoking methods or path expressions on objects may be costly to execute, traditional heuristics for optimizing recursive queries are no longer valid. In this paper we propose a cost-based optimization method which handles object-oriented recursive queries. In particular, it is ...

13 Query execution in prism and seaview: a cost analysis

Brajendra Panda, William Perrizo

February 1995 Proceedings of the 1995 ACM symposium on Applied computing

Full text available: pdf(699.80 KB) Additional Information: full citation, references, index terms

Keywords: data and user classification, military securitY policies, multilevel databases, query processing

14 Web services and performance evaluation: Indexing web access-logs for pattern queries

Alexandros Nanopoulos, Yannis Manolopoulos, Maciej Zakrzewicz, Tadeusz Morzy November 2002 Proceedings of the fourth international workshop on Web information and data management

Full text available: pdf(187.24 KB) Additional Information: full citation, abstract, references, index terms

In this paper, we develop a new indexing method for large web access-logs. We are concerned with pattern queries, which advocate the search for access sequences that contain certain query patterns. This kind of queries find applications in processing web-log mining results (e.g., finding typical/atypical access-sequences). The proposed method focuses on scalability to web-logs' sizes. For this reason, we examine the gains due to signature-trees, which can further improve the scalability to very ...

15 <u>Cost-based query optimization for metadata repositories</u>
Martin Staudt, René Soiron, Christoph Quix, Matthias Jarke
March 1999 **ACM SIGAPP Applied Computing Review**, Volume 7 Issue 2

Full text available: pdf(1.04 MB)

Additional Information: full citation, abstract, index terms

Query optimization strategies for repository systems must take into account the rich and often unpredictable structure of metadata, as well as supporting complex analysis of relationships between those structures. This paper describes rationale, design, and system integration of a cost-based query optimizer offered in ConceptBase, a metadata manager that supports these capabilities by a deductive object-oriented data model. In contrast to most implemented DBMS, the optimizer is not based on the ...

16 <u>Database principles: A mapping mechanism to support bitmap index and other auxiliary structures on tables stored as primary B*-trees</u>

Eugene Inseok Chong, Chuck Freiwald, Anh-Tuan Tran, Jagannathan Srinivasan, Aravind Yalamanchi, Ramkumar Krishnan, Souripriya Das, Mahesh Jagannath, Richard Jiang June 2003 **ACM SIGMOD Record**, Volume 32 Issue 2

Full text available: pdf(198.55 KB) Additional Information: full citation, abstract, references

Any auxiliary structure, such as a bitmap or a B⁺-tree index, that refers to rows of a table stored as a primary B⁺-tree (e.g., tables with clustered index in Microsoft SQL Server, or index-organized tables in Oracle) by their physical addresses would require updates due to inherent volatility of those addresses. To address this problem, we propose a mapping mechanism that 1) introduces a single mapping table, with each row holding one key value from the prima ...

17 Query optimization in a memory-resident domain relational calculus database system Kyu-Young Whang, Ravi Krishnamurthy



Full text available: pdf(2.46 MB)

Additional Information: <u>full citation</u>, <u>abstract</u>, <u>references</u>, <u>citings</u>, <u>index</u> terms

We present techniques for optimizing queries in memory-resident database systems. Optimization techniques in memory-resident database systems differ significantly from those in conventional disk-resident database systems. In this paper we address the following aspects of query optimization in such systems and present specific solutions for them: (1) a new approach to developing a CPU-intensive cost model; (2) new optimization strategies for main-memory query processing; (3) new insight into ...

18 Optimal indexing using near-minimal space

C. Heeren, H. V. Jagadish, L. Pitt

June 2003 Proceedings of the twenty-second ACM SIGMOD-SIGACT-SIGART symposium on Principles of database systems

Full text available: pdf(208.43 KB) Additional Information: full citation, abstract, references, index terms

We consider the index selection problem. Given either a fixed query workload or an unknown probability distribution on possible future queries, and a bound ${\it B}$ on how much space is available to build indices, we seek to build a collection of indices for which the average query response time is minimized. We give strong negative and positive peformance bounds. Let ${\it m}$ be the number of queries in the workload. We show how to obtain with high probability a collection of indices using space ...

19 Parametric query optimization

Yannis E. Ioannidis, Raymond T. Ng, Kyuseok Shim, Timos K. Sellis

May 1997 The VLDB Journal — The International Journal on Very Large Data Bases, Volume 6 Issue 2

Full text available: pdf(378.53 KB) Additional Information: full citation, abstract, index terms

In most database systems, the values of many important run-time parameters of the system, the data, or the query are unknown at query optimization time. Parametric query

optimization attempts to identify at compile time several execution plans, each one of which is optimal for a subset of all possible values of the run-time parameters. The goal is that at run time, when the actual parameter values are known, the appropriate plan should be identifiable with essentially no overhead. We present a g ...

20 Query processing in a multimedia document system

Elisa Bertino, Fausto Rabbiti, Simon Gibbs

January 1988 ACM Transactions on Information Systems (TOIS), Volume 6 Issue 1

Full text available: pdf(2.94 MB)

Additional Information: full citation, abstract, references, citings, index terms, review

Query processing in a multimedia document system is described. Multimedia documents are information objects containing formatted data, text, image, graphics, and voice. The query language is based on a conceptual document model that allows the users to formulate queries on both document content and structure. The architecture of the system is outlined, with focus on the storage organization in which both optical and magnetic devices can coexist. Query processing and the different strategies ...

Results 1 - 20 of 200

Result page: **1** <u>2</u> <u>3</u> <u>4</u> <u>5</u> <u>6</u> <u>7</u> <u>8</u> <u>9</u> <u>10</u> next

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